Norm equations and More in Oscar

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Claus Fieker [Norm equations and More in Oscar](#page-47-0)

Topics:

- **o** Oscar
- Norm Equations
- ... and more

[Norm Equations](#page-7-0) [Diophantine Norm Equations](#page-40-0) [More on Oscar](#page-46-0)

Develop a visionary, next generation, open source computer algebra system, integrating all systems, libraries and packages developed within the TRR. (Be able to compete with Magma in our area of expertise.)

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What is Oscar?

<http://oscar-system.org/><https://oscar-system.github.io/oscar-website/>

- (software) project of the CRC 195
- **•** funded by DFG
- in Julia
- \bullet funding (planned) 2017 2028
- in three phases

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Oscar

```
ju lia > Pkg. add ("Oscar")
\ldots [wait some time] \ldotsi u lia > using Oscar
```

```
\vert \ \ \vert \ \ \vert \ \ \vert \ \ \rangle --- \ \setminus \ \ \vert \ \ \ \ \ \ \ \ \ \ \vert - \ \ \vert \ \ \ \vert \ \ \ \ \vert Type "?Oscar" for more information
```
/ \ / | / | / \ | \ | Combining ANTIC , GAP, Polymake , S i n g u l a r | |_| | ___) | |___ / ___ \| _ < | Manual: https://docs.oscar—system.org \ / | / \ / / \ \ | \ \ | 1.2.0 −DEV #ma s t e r f 4 8 9 2 2 0 2024−06−24

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What is Oscar?

O pen

S ource

- C omputer
- A lgebra
- R esearch

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Why Julia?

- **o** Interactive
- As fast as C
- Solves 2-language problem
- Not maintained by us
- **o** Modern
- Interoperates well with C
- (Originates at the MIT, group of Alan Edelman)

Intro

.

Given some finite extension

 A/B

with a norm map $N : A \to B$ and some $b \in B$, find one/all $a \in A$ s.th.

 $N(a) = b$

First Examples

- K/\mathbb{Q} a number field (absolute norm equation)
- \bullet K/k number fields (relative norm equation)
- \bullet \mathcal{O}_K/\mathbb{Z} (Diophantine case)
- \bullet $\mathbb{Z}[\alpha]/\mathbb{Z}$ (for Thue equations)
- K/k finite fields
- \bullet K/k local fields
- $\mathbb{I}_K/\mathbb{I}_k$ idele-ic case

... and the decision problems as well. ... and function fields.

Motivation/Applications

- In representation theory of finite groups: norm equations determine/ are used to find minimal fields for the representation.
- \bullet Points on conics (= isotropic vectors for quad. forms) are used in elliptic curves.
- Solving Thue equations starts by solving norm equations.
- (Some) Embedding problems reduce to norm equations
- Base case of Galois cohomology

Broadly

Different cases require different techniques and show different runtime characteristica:

- K/k number fields, K/k normal: S-units and linear algebra
- \bullet K/k number fields, not normal: very interesting!
- \bullet \mathcal{O}_K/\mathbb{Z} : S-units and lattice points in polytopes.
- \bullet K/k finite fields: irreducibility and roots.
- \bullet K/k local fields: finite fields and linear algebra.
- \bullet $\mathbb{I}_{K}/\mathbb{I}_{k}$: local fields and their mult. group.

Decision Problems

- **•** finite fields: no problem, norm is surjective
- \bullet K/k number fields, cyclic case: Hasse norm theorem: solvable iff locally solvable everywhere
- \bullet K/k normal: "known" obstacle to Hasse norm theorem, the Knot
- \bullet K/k not normal: Knot is known in some cases.
- \bullet K/k local, unramified: valuation only, trivial

Finite Fields

 K/k finite fields, $b \in k$, $b \neq 0$.

- pick random irreducible monic polynomial f of degree $[K : k]$ with constant term $(-1)^{[K:k]}b$
- \bullet return a root a of f
- Steel (approx. 2002) from Shoup (s.t. before), folklore

In almost all other cases we have to work.

Number Fields and Rings

Now fix a finite extension K/k of number fields. Fundamental idea: $a \in K$ s.th. $N(a) = b$ is a S-unit. Find a suitable set S, compute the S-unit group and find a . Problem:

- \bullet S will depend on b, but may depend on K as well
- for number rings, we also want integrality
- \bullet for non Dedekind domains using ideals (and S -units) is tricky

Roughly

 T a set of ideals in K and S in k s.th.

```
N \cdot T-units in K \rightarrow S-units in k
```
is well defined. Assume that the S - and T-unit groups are given algorithmically, ie.

- as an abstract abelian group
- with disc. log and disc. exponential

Then N can be constructed explicitly as a map between abelian groups, and the rest is easy.

Example: Number Field

Want

$$
N(b) = 31
$$

for $b \in \mathbb{Q}(\sqrt{2})$ 10):

```
julia> k, a = quadratic_field(10)
(Real quadratic field defined by x^2 - 10, sqrt(10))
```

```
julia> zk = maximal-order(k)Maximal order of Real quadratic field defined by x^2 - 10
with basis AbsSimpleNumFieldElem[1, sqrt(10)]
```
julia $> a = 31$

Example: Number Field

```
We choose S = \{31\} and T the primes above.
julia> S, mS = sunit_group([31])
(Z/2 x Z, SUnits map of Rational field for ZZRingElem[31]
)
```

```
julia> T, mT = sunit_group(prime_ideals_over(zk, 31))
(Z/2 x Z^(3), SUnits map of k for AbsSimpleNumFieldOrderIdeal[<31, sqrt(1
Norm: 31
Minimum: 31
two normal wrt: 31, <31, sqrt(10) + 14>
Norm: 31
Minimum: 31
two normal wrt: 31]
)
```
Example: Number Field

Now we set up the norm map: Using mT to map abstract generators for the T -unit group into elements on K , then applying the norm and finally using the disc. log in the S-units (of \mathbb{O}), via the preimage of mS. All of this is collected in an (abstract) homomorphism.

```
julia> N = hom(T, S, [preimage(mS, norm(mT(x))))for x = \text{gens}(T))
Map
```
from $7/2 \times 7^{\circ}(3)$ to Z/2 x Z

Example: Number Field

Now testing for solvability just requires to write a as an element of S and using the abstract norm map... In order to get a solution, the abstract preimage has to be converted into an element of K again.

```
julia> preimage(mS, a)
Abelian group element [0, 1]
```

```
julia> has_preimage(N, ans)
(true, Abelian group element [0, -1, 0, 1])
```

```
julia> mT(ans[2])
3*sqrt(10) + 11
```

```
julia> norm(ans)
31
```
Problem(s)

Trying for $K = \mathbb{Q}(\sqrt{2})$ $34)$, we see that

$$
\mathcal{O}_K^* = \langle -1, 35 - 6\sqrt{34} \rangle
$$

Since $N(35-6)$ √ $34)=1$, there is no integral element with norm $-1.$ However

$$
N(\frac{1}{3}\sqrt{34}-\frac{5}{3})=-1
$$

so in general finding S and T is non-trivial, T cannot just depend on the RHS. Furthermore, since units can be **extremely** large, solutions can not be expected to be small.

Factored elements

One (well known) improvement is to use factored elements: since units are (frequently) huge, we use a multiplicative representation

$\prod \alpha_i^{n_i}$

for (smallish) α_i and (largeish) exponents $n_i \in \mathbb{Z}$.

- Magma: ProductRepresentation or Raw
- Pari/gp: $\mathbb{Z}[K]$ presentation
- Oscar: factorised elements, FacElem

We need to replace sunit group by sunit group fac elem...or even sunit mod_units_group_fac_elem. And add evaluate at the end. Note: we can also obtain a compact presentation.

Norm Equations - Factored Elements

```
julia> T, mT = sunit_group_fac_elem(prime_ideals_over(maximal_order(k), 31))
(Z/2 \times Z^*(3), SUnits (in factored form) map of Factored elements over Real
)
...
julia> N = hom(T, S, [preimage(mS, norm(mT(t))) for t = genes(T)])...
julia> preimage(N, ans)
Abelian group element [0, -1, 0, 1]julia> mT(ans)
(2*sqrt(10))^2*(sqrt(10) + 17)^{-1*7-1*7-1*67-1*67-1*4-2*13^3*(sqrt(10) + 9))
```

```
julia> evaluate(ans)
3*sqrt(10) + 11
```
Ideals

So:

 $N(a) = b$

for $a \neq T$ -unit. How do we find T ? Observation:

- \bullet wlog b is integral
- if $N(a) = b$, then $N(a\mathcal{O}_K) = b\mathcal{O}_k$ as well
- on ideals we have unique factorisation

So, assume we have a solution $a \in K$ and any integral ideal $\mathfrak A$ of the correct norm. Then $N(a\mathfrak{A}^{-1})=\mathcal{O}_k$.

Ideals - Hilbert 90 and Normal Fields

Simplifying to $k = \mathbb{Q}$ and K/\mathbb{Q} normal, we have Hilbert 90 for ideals:

Theorem $N(\mathfrak{A})=1$ iff $\mathfrak{A}=\prod P_i^{1-\sigma_i}$.

Or even more fancy $(I_K$ the group of invertable ideals of K):

Theorem

Both $H^1(\mathsf{Gal}(K/k),I_K)$ and $H_1(\mathsf{Gal}(K/k),I_K)$ are trivial.

Ideals - Hilbert 90 and Normal Fields

So
$$
N(a) = b
$$
, $N(\mathfrak{A}) = b\mathcal{O}_k$ so $N(a\mathfrak{A}^{-1}) = \mathcal{O}_k$ and

$$
a\mathfrak{A}^{-1} = \prod P_i^{1-\sigma_i}
$$

If $\mathsf{CI}_K = \langle \mathfrak{B}_j | 1 \leq j \leq k \rangle$, then for any P_i there is some α_i s.th. $P_i = \alpha_i \prod \mathfrak{B}_j^{n_{i,j}}$ $_j^{n_{i,j}}$, thus

$$
a\mathfrak{A}^{-1} = \prod \alpha_i^{1-\sigma_i} \prod \mathfrak{B}_i^{n_{i,j}(1-\sigma_i)}
$$

Sorting:

$$
a \prod \alpha_i^{\sigma_i - 1} \mathcal{O}_K = \mathfrak{A} \prod \mathfrak{B}_i^{n_{i,j}(1 - \sigma_i)}
$$

The LHS is now a solution as well and the RHS has a known support.

Hence: we have a set T of prime ideals s.th. if there is a solution, there is one with support in $T!$

 T depends on the RHS b and the class group.

Non-normal Fields

- Siegel (1973), Bartel (1980): choose T as the set of primes of norm bounded by the Minkowski constant of the normal closure.
- \bullet Simon (1998) choose T to generate the class groups of all relative cyclic subfields of the normal closure.

Pros and Cons:

- Bartel: no need to compute in a large field, many primes
- Simon: many class groups, small set T , can use GRH

Would like small set T , use GRH and do computations in K only.

Norm 1 Ideals

Fix K/k finite. $N = N_{K/k}$ denotes the norm (on elements and ideals). For ideals $\mathfrak A$ and $\mathfrak B$ of norm $1 = \mathcal O_k$ s.th. $|\mathfrak A| = |\mathfrak B|$ in the class group we get

 $\mathfrak{A} = \beta \mathfrak{B}$

and $N(\beta)\mathcal{O}_k = \mathcal{O}_k$, so $N(\beta) \in U_k$, a unit:

Norm 1 Ideals

Collapsing from the left:

And since $U_k^n < N(U_K)$, we get U_k $\frac{U_{k}}{N(U_{K})}$ is finite ${a|N(a) \in U_k}$ $\frac{\lbrace a \rbrace N(a) \in U_k \rbrace}{U_K \lbrace a \rbrace N(a) = 1 \rbrace}$ is finite \bullet X, the subgroup of Cl_K gen. by ideals of norm 1 is finite,

Norm 1 Ideals

Yields:

$$
\mathsf{Cl}_K^1 := \frac{\{\mathfrak{A} \mid N(\mathfrak{A}) = \mathcal{O}_k\}}{\{a \mid N(a) = 1\}}
$$

is finite!

Partly constructive, $X = \{x \in \mathsf{Cl}_K \mid \exists \mathfrak{A} \in x, N(\mathfrak{A}) = 1\}$

- we can generate ideals of norm 1
- given two such ideals we can check equality in X (and $\mathsf{Cl}_K^1)$

So we can just do this to get a subgroup, but completeness? Similarly, we can test if $\mathfrak A$ and $\mathfrak B$ coincide in Cl $^1_K\colon \mathfrak A=\beta\mathfrak B$ and $N(\beta)\in N(U_K).$ So we can obtain a subgroup of Cl 1_K , but completeness?

Assume Cl $^1_K=\langle\mathfrak{A}_i|i\rangle$, then we solve norm equations as before: <code>(wlog</code> $b\in\mathcal{O}_k$), assume $N(a) = b$

- **1** find $\mathfrak{A} \leq \mathcal{O}_K$ s.th. $N(\mathfrak{A}) = b\mathcal{O}_k$
- 2 then $N(a\mathfrak{A}^{-1})=\mathcal{O}_k$
- $\textbf{3}$ thus $a\mathfrak{A}^{-1}\in \textsf{Cl}_K^1$
- $\bullet \ \ \hbox{there is} \ \beta, \ N(\beta) = 1 \ \hbox{and} \ a \mathfrak A^{-1} = \beta \mathfrak A_{i}$
- **5** then $a\beta^{-1}\mathcal{O}_K = \mathfrak{A}\mathfrak{A}_i$ yields a solution with a known support.

The set S

Then $P\mathcal{O}_\Gamma=\prod Q_i.$ G operates transitively on the primes, so for $Q=Q_1$ we have $Q_i = Q_1^{s_i}$ for $s_i \in G$, hence, $\sigma := \sum s_i \in \mathbb{Z}[G]$:

$$
P\mathcal{O}_\Gamma=Q^\sigma
$$

Let $\mathfrak A$ be an ideal in K, $N(\mathfrak A) = \mathcal O_k$ and supported only at primes above p.

$$
\mathfrak{A}\mathcal{O}_{\Gamma} = \prod P_i^{n_i} \mathcal{O}_{\Gamma} = \prod Q^{n_i \sigma_i} = Q^{\sum n_i \sigma_i} =: Q^{\tau}
$$

Since $N(\mathfrak{A}) = \mathcal{O}_k$, also $N_{\Gamma/k}(\mathfrak{A} \mathcal{O}_{\Gamma}) = \mathcal{O}_k$ as well. So

$$
\tau \in I_G = \langle 1-s | s \in G \rangle \leq \mathbb{Z}[G],
$$

the augmentation ideal.

 $\mathfrak{A} \subset K$, so τ is stable under Aut (Γ/K) , $\tau s = \tau$ for all $s \in \text{Aut}(\Gamma/K)$! Thus: $\alpha^{\tau} \in K$ and $N(\alpha^{\tau}) = 1$ for all $\alpha \in \Gamma$.

Let Cl $_{\Gamma}=\langle [S_i]|i\rangle$ for unramified ideals S_i and $\mathfrak{A}\leq\mathcal{O}_K$ of norm 1 as above, so $\exists S = \prod S_i^{n_i}$ and $\alpha \in \Gamma$: $\mathfrak{A} \mathcal{O}_\Gamma = R^\tau = (\alpha S)^\tau = \alpha^\tau S^\tau$

Thus in Cl 1_K all (unramified) ideals of norm 1 come from generators of the class group of Γ. — GRH or unconditional.

(The (few) (fintely many) ramified ideals of norm 1 are easily added.)

Let $\mathfrak m$ be an integral ideal in k s.th. for all units $u\in\mathcal O^*_k$, $u\equiv 1\bmod \mathfrak m$ we have that $u = v^n$ for $n = [K : k]$ Let $X \n\leq C |_{mQ_V}$ be subgroup of rays containing ideals of norm 1. Then

$$
1\to X\to {\sf Cl}_{\mathfrak m\mathcal O_K}
$$

is exact:

 $\mathfrak{A} = \mathfrak{B}$ in $\mathsf{Cl}_{\mathfrak{m}\mathcal{O}_K}$ implies $\mathfrak{A} = \beta \mathfrak{B}$ and $\beta = 1 \mod \mathfrak{m}\mathcal{O}_K$. $N(\mathfrak{A}) = N(\mathfrak{B}) = \mathcal{O}_k$ implies $N(\beta) \in U_k$. Since $N(\beta) = 1 \bmod{\mathfrak{m}}$. So $N(\beta) = \epsilon^n$ and $N(\beta/\epsilon) = 1$ and $\mathfrak{A}=\mathfrak{B}\in\mathsf{Cl}_K^1$ as well. This is easier to work with than Cl_K^1 directly - but misses the primes in $\mathfrak{m}.$

The algorithm(s)

Solve $N(a) = b$.

1 Find a suitable m

$$
\bullet \ \ S = \{\}, \ X = \langle \mathcal{O}_K \rangle \leq \mathsf{Cl}_{\mathfrak{m}}
$$

 \bullet for p (unramified) primes in k (coprime to m) do

\n- **①**
$$
p\mathcal{O}_K = \prod P^i
$$
 with $N(P_i) = p^{f_i}$
\n- **②** Let $n_{i,j}$ a Z-basis for $\sum n_{i,j} f_i = 0$
\n- **③** if $\prod P_i^{n_{i,j}} \notin X$, then add p to S and enlarge X
\n

We can

- use the Minkowski/ Bach/ Belabas et. al./ \ldots bound for S ,
- stop the search when X did not change for ? primes p.

We have to supplement with the primes in m and the ramified ones. T is the set of primes above S .

Knots - the Decision Problem

Let \mathbb{I}_{K} , resp. \mathbb{I}_{k} the idele groups, then Scholz (1936) defined the (number) knot

 $\delta_{K/k} := N(\mathbb{I}_K)/N(K^*)$

to measure the error in the Hasse norm theorem:

Theorem

For cyclic extensions K/k , the knot is trivial, hence solvability can be tested locally.

Well, actually, not, he studied:

Wir nennen K_0 die Restklassengruppe der Normreste nach den Zahlnormen den (Gesamt-, Zahl-) Knoten $K = K_{\alpha}$ von K.

(We call the quotient group of norm residues modulo norms the (total-, number-) knot.) Free of ideles.

Knots - the Decision Problem

$$
\delta_{K/k} := N(\mathbb{I}_K)/N(K^*)
$$

- The knot is trivial for cyclic extensions.
- There are infinitely many bi-quadratic fields where the knot is non trivial, but also infinitely many where the norm theorem holds. Newton et. al. studied this quantitatively.
- Testing local solvability is also not always easy.
- Jehne defined many more knots in 1979.

For K/k normal, the knot can be computed classically (Tate)

$$
1 \to \delta_{K/k} \to H^2(G, \mathbb{Q}/\mathbb{Z}) \to \oplus H^2(G_p, \mathbb{Q}/\mathbb{Z})
$$

(The sum runs over all places of K and G_p are the decomposition groups, the local Galois groups.) For abelian G , this is easy, for general G one needs more group theory (Schur multipliers, group cohomology).

Strategy

If not locally solvable: return fail. Start with choosing S to contain

- all primes dividing the RHS
- all ramified primes
- enough primes to likely generate Cl_K^1

And try to find a solution using U_S units. If this fails then

- \bullet if knot is trivial: increase S until it works
- \bullet use GRH (or not) to enlarge S and try again

Integral Norm Equations

Let K/k be finite and $b \in \mathcal{O}_k$. Find

 $N(a) = b$

for $a \in \mathcal{O}_K$. Modulo units in K, this equation has only fin. many solutions. Classical: turn into lattice problem. Let the units operate, to obtain a finite search domain.

Solve by enumeration:

- bound into a box (classical)
- cover by ellipsoids (Fincke/Pohst (1988), Jurk (1993), F.(1997)).

Very nice, beautiful pictures, slow.

Using Class Group

 $N(a) = b$

For $a \in \mathcal{O}_K$. Any solution a generates a principal ideal $a\mathcal{O}_K$ of norm $b\mathcal{O}_k$. So

- **1** find all integral ideals of norm bO_k
- 2 for each ideal test if principal
- **3** for each principal ideal test if the generator can be made to work

Let P_i be the primes in \mathcal{O}_K dividing a prime in \mathcal{O}_k of the support of b. We have $N(P_i)=p_{j_i}^{f_i}$ $j_i^{f_i}$ and, $P_i \rightsquigarrow c_i \in \mathsf{Cl}_K$. To list all integral principal ideals of the correct norm is a classical combinatorial

problem:

Find $n = (n_i)_i$ s.th.

• $n_i > 0$ for all i (integrality)

• Let
$$
A = (v_{p_j}(P_i))_{i,j}
$$
 and $An = (v_{p_j}(b))_j$ (norm)

 $\sum n_i c_i = 0$ (principality)

Then $\prod P_i^{n_i}$ is an integral principal ideal of the correct norm (up to units).

Using S-Units

The same, without the class group, using S -units directly:

$$
\bullet \text{ find } S := \{P \leq \mathcal{O}_K | v_P(b) > 0\}
$$

compute $U_S/U_K = \langle \epsilon_i \mid i \rangle$

$$
\bullet\ A:=(v_{P_j}(\epsilon_i))_{i,j},\ B:=(v_p(N(\epsilon_j)))_{i,j}
$$

• solve $An \geq 0$ s.th. $Bn = (v_p(b))_p$ and (try to) adjust by units

In all cases, this is a combinatorial problem: points in a lattice (in ellipsoids) or points in a polytope.

Non-maximal Order

Final case:

$$
N(a) = b
$$

for $a \in \mathcal{O}$ a non-maximal order. This is used e.g. in Thue equations where $\mathcal{O} = \mathbb{Z}[\alpha]$ is an equation order.

If $N(a) = b$ for $a \in \mathcal{O} \subseteq \mathcal{O}_K$, then any solution is also one from the last case. Step 1: solve in the maximal order, find all solutions. Since $U_K/{\mathcal O}^*$ is finite, any solution in ${\mathcal O}$ is obtained from one in ${\mathcal O}_K$ and a unit $\epsilon \in U_K/{\mathcal O}^*$ of norm 1.

For the final step, assuming b is coprime to the conductor f of O in \mathcal{O}_K , then a is also coprime, so

$$
a \in (\mathcal{O}_K/\mathfrak{f})^*
$$

$$
\downarrow
$$

$$
0 \longrightarrow \mathcal{O}^* \longrightarrow U_K \longrightarrow (\mathcal{O}_K/\mathfrak{f})^*/(\mathcal{O}/\mathfrak{f})^* =: X
$$

This is always used to compute \mathcal{O}^* or $\mathcal{O}_K^*/\mathcal{O}^*$, but can also be used in the last step: solutions in O correspond to preimages in U_K of $a \in X$

- https://oscar-system.org
- The Oscar book: The Computer Algebra System OSCAR
- development on https://github.com/oscar-system/Oscar.jl
- most number theory https://github.com/thofma/Hecke.jl
- \bullet foundations on https://github.com/Nemocas/AbstractAlgebra.jl
- optimizations https://github.com/Nemocas/Nemo.jl
- active slack community https://oscar-system.org/slack

Features (Number Theory)

- number fields, (non)-simple extensions
- orders, maximal order
- class groups, $(S₋)$ units
- Galois groups, automorphisms
- constructive class field theory
- localizations and completions
- **•** Galois cohomology
- algebras and orders

...

- lattices, automorphisms, isomorphisms
- **o** function fields and orders
- verified real computations using arb