

# Efficient Hyperelliptic Arithmetic Using Balanced Representation for Divisors

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# Overview

- 1 Introduction
- 2 Addition
- 3 Comparison

## Plane Models

We will consider a genus  $g$  hyperelliptic curve  $C$  defined over a field  $k$  with  $\text{char}(k) \neq 2$ . We can assume that  $C$  is given by a plane model

$$C : y^2 = F(x),$$

where  $F$  is a polynomial in  $k[x]$  with no repeated roots.

If  $P = (x, y)$  is a point on the curve, then  $\bar{P} = (x, -y)$  also lies on the curve and is called hyperelliptic conjugate of  $P$ .

## Taxonomy

- If  $\deg(F)$  is  $2g + 1$ , this is an imaginary model.  $C$  will have 1 point at infinity.
- If  $\deg(F)$  is  $2g + 2$ , this is a real model.  $C$  will have 2 points at infinity.

## The divisor class group

- The *group of divisors* on  $C$  is the group of finite formal sums  $D = \sum n_i P_i$ , for integers  $n_i$  and points  $P_i$  on  $C(\bar{k})$ .  
 $\deg(D) = \sum n_i$ .
- To every rational function  $f$  in  $C(\bar{k})^*$ , one can associate a divisor

$$\operatorname{div}(f) = \sum_{P \in C(\bar{k})} \operatorname{ord}_P(f).$$

The set of divisors associated to all the functions in  $C(\bar{k})^*$  forms the subgroup of principal divisors.

- The divisor class group of  $C$  is the quotient group of the group of divisors modulo the subgroup of principal divisors. The class of  $D$  will be denoted  $[D]$ .

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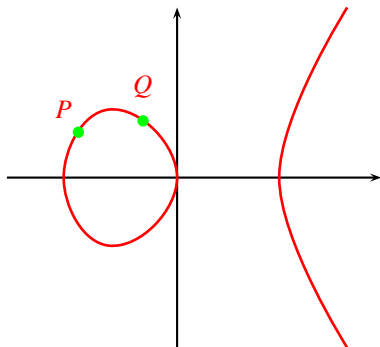
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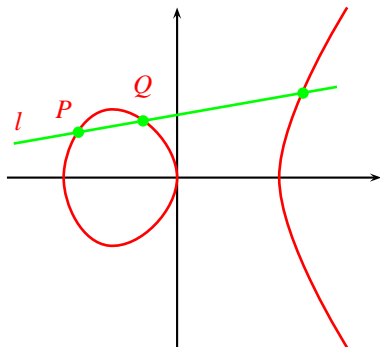
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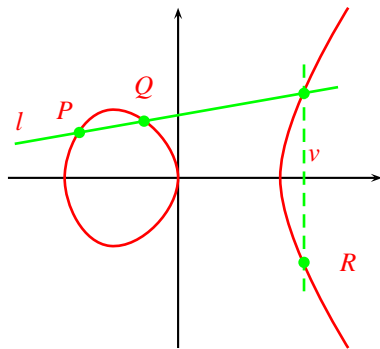
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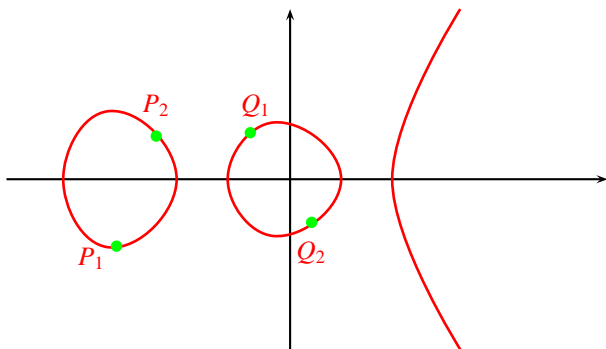
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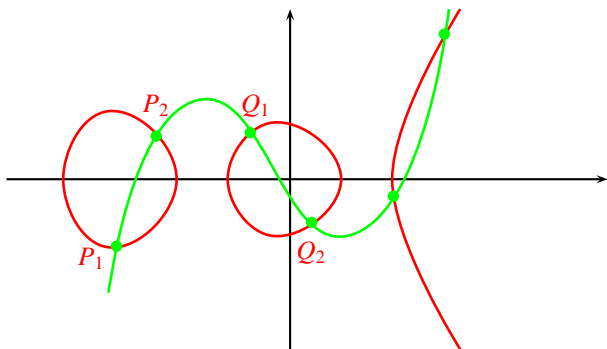
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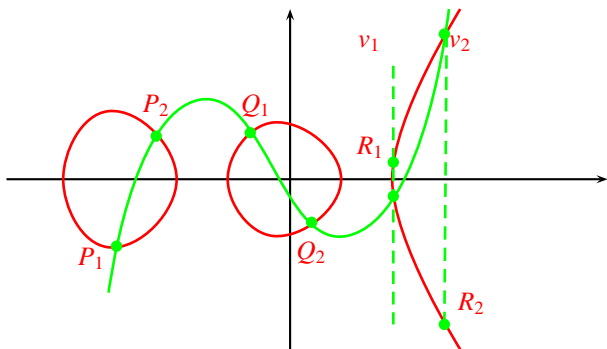
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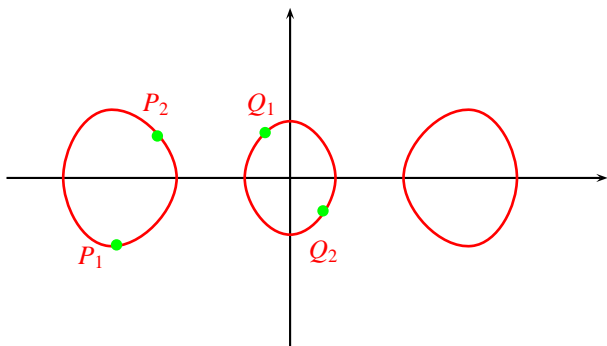
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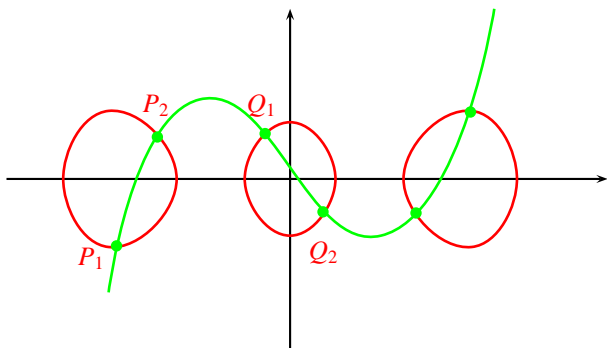
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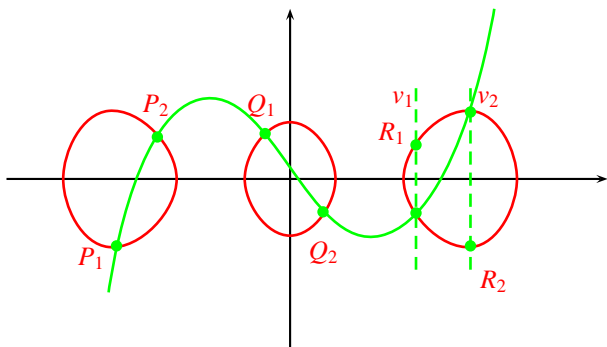
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## Definition

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## Theorem

Let  $D_\infty$  be a  $k$ -rational degree  $g$  divisor, and let  $D \in \text{Div}^0(C)$  be a  $k$ -rational divisor on the hyperelliptic curve  $C$ . Then  $[D]$  has a unique representative in  $\text{Cl}^0(C)$  of the form  $[D_0 - D_\infty]$ , where  $D_0$  is an effective  $k$ -rational divisor of degree  $g$  whose affine part is reduced.

## The base divisor

- If  $C$  is given by an imaginary model, then  $D_\infty = g\infty$ .
- If  $C$  is given by a real model denote its points at infinity as  $\infty^+$  and  $\infty^-$ . Then
  - $D_\infty = \frac{g}{2}(\infty^+ + \infty^-)$  if  $g$  is even.
  - $D_\infty = \frac{g+1}{2}\infty^+ + \frac{g-1}{2}\infty^-$  if  $g$  is odd.

## Mumford's Representation

To every pair of polynomials  $(u(x), v(x))$  such that

$$u(x) \text{ divides } F(x) - v(x)^2, \quad (1)$$

we associate a divisor as follows

$$\text{If } u(x) = \prod_i (x - r_i), \text{ then } (u(x), v(x)) \mapsto \sum_i (r_i, v(r_i)).$$

We say that the polynomials  $(u, v)$  are the Mumford representation of  $D$ , and denote this as  $D = \text{div}(u, v)$ . Every affine semi-reduced divisor has a Mumford representation.

## Problem

Fix a degree  $g$  divisor  $D_\infty$ . Given two effective degree  $g$  divisors with reduced affine part  $D_1$  and  $D_2$ , find an effective degree  $g$  divisor  $D_3$  with reduced affine part such that

$$[D_1 - D_\infty] + [D_2 - D_\infty] = [D_3 - D_\infty].$$

## Equivalently

To add the divisor classes  $[D_1 - D_\infty]$  and  $[D_2 - D_\infty]$ , one calculates  $D_3$  satisfying

$$[D_1 + D_2] = [D_3 + D_\infty].$$

- 1 Given the Mumford representation of  $D_1$  and  $D_2$ , find the Mumford representation of  $D_1 + D_2$ .
- 2 From the Mumford representation of  $D_1 + D_2$ , find the appropriate  $D_3$ . This is done using the reduction algorithms.

## Reduction

Let  $D_0 = \text{div}(u_0, v_0)$  be a divisor of degree  $d_0 \geq g + 2$  ( $d_0 \geq g + 1$ ). By definition of the Mumford representation, the divisor of  $y - v_0(x)$  has (generically) the form

$$\text{div}(y - v_0(x)) = D_0 + D_1 - \frac{d_0 + d_1}{2}(\infty^+ + \infty^-),$$

where  $D_1$  is an affine semi-reduced divisor. This implies

$$[D_0] = [\overline{D}_1 + \frac{d_0 - d_1}{2}(\infty^+ + \infty^-)].$$

The affine zeros of  $y - v_0(x)$  are found solving  $v_0(x)^2 - F(x) = 0$ .

## Composition and reduction

- If the divisor  $D = \operatorname{div}(u, v)$  has degree at most  $g + 1$ , then reduction using  $y - v(x)$  does not work.
- For instance if  $D$  has degree  $g + 1$ , then  $\deg(v) \leq g$ , so  $v^2 - F$  will have  $2g + 2$  affine zeros, and we get another divisor of degree  $g + 1$ .
- We have cancelation in  $v^2 - F$  if and only if the leading term of  $p$  is  $F_{2g+2}^{1/2} x^{g+1}$ .
- The function  $y - p(x)$  has different order at  $\infty^+$  and  $\infty^-$  if and only if the leading term of  $p$  is  $F_{2g+2}^{1/2} x^{g+1}$ .

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## This motivates

Let  $H^+(x)$  be the polynomial with leading term  $F_{2g+2}^{1/2}x^{g+1}$  such that  $(H^+)^2 - F$  has minimal degree. Given  $D = \text{div}(u, v)$  of degree at most  $g + 1$ , use the polynomial

$$p(x) = H^+ + (v - H^+ \pmod{u})$$

to perform a reduction (This is a  $\text{red}_\infty$  step).

## Generically

- If  $D_0$  has degree  $g$ , then typically  $[D_0] = [D_1 + (\infty^+ - \infty^-)]$ .
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## Generical addition for even genus

- Given  $[D_1 - D_\infty]$  and  $[D_2 - D_\infty]$  find the Mumford representation of  $D_1 + D_2$ .
- Reduce until  $[D_1 + D_2] = [D_3 + (g/2)(\infty^+ + \infty^-)]$ .
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- Reduce until  $[D_1 + D_2] = [D'_3 + (g - 1)/2(\infty^+ + \infty^-)]$ .
- Use composition-and-reduction to get  $[D'_3] = [D_3 + \infty^+]$ .
- Now  $[D_1 + D_2] = [D'_3 + (g + 1)/2\infty^+ + (g - 1)/2\infty^-]$ .
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## In the olden days

Previous authors used the base divisor  $D_\infty = g_\infty^+$  instead of the “balanced” divisor we proposed. We will show that this is slower than our approach.

## In the even genus case

Using a a balanced representation, we finished when we found  $D_4$  such that

$$D_1 + D_2 \equiv D_4 + \frac{g}{2}(\infty^+ + \infty^-).$$

If we wanted to use  $D_\infty = g\infty^+$  instead, we'd get

$$[D_1 - D_\infty] + [D_2 - D_\infty] = [D_4 - D_\infty] + \frac{g}{2}(\infty^- - \infty^+),$$

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## Divisor inversion

It is important to be able to invert elements in a group (window methods, signed representations, etc).

Using our representation, this can be done with 0 or 1  $\text{red}_\infty$  steps.

- If the genus is even, then  $D_\infty = \overline{D_\infty}$ .
- If the genus is odd, then  $D_\infty = \overline{D_\infty} + (\infty^+ - \infty^-)$ .

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- If the genus is odd, then  $D_\infty = \overline{D_\infty} + (\infty^+ - \infty^-)$ .

Using  $g\infty^+$  as base divisor, it takes  $g$  applications of  $\text{red}_\infty$ .

- If  $D_\infty = g\infty^+$ , then  $D_\infty = \overline{D_\infty} + g(\infty^+ - \infty^-)$ .

## Comparison

In a genus 2 curve if  $S = M$  and  $I = 4M$  then balanced representations give a saving of around 15% for addition and 13% for doubling (if  $I = 30M$  the savings become 62% and 58% respectively).

	Imaginary	Balanced	Non-balanced
Addition	1I, 2S, 22M	1I, 2S, 26M	2I, 4S, 30M
Doubling	1I, 5S, 22M	1I, 4S, 28M	2I, 6S, 32M
Inversion	0	0	2I, 4S, 8M

**Table:** Operation counts for genus 2 arithmetic.

## Remarks

- The analogue of a baby-step in the imaginary model is addition of  $P - \infty$ .
- If the genus is even and the points at infinity are not rational, baby-steps are not necessary.
- Implemented in Magma V2.12, July 2005.
- One can efficiently implement pairings on hyperelliptic curves given by a real model (upcoming article with S. Galbraith and X. Lin ).

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# Questions?

Thank you for your attention